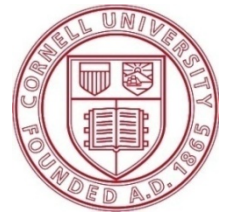
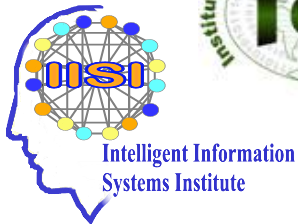


EMERGENCE OF INTELLIGENT MACHINES: CHALLENGES AND OPPORTUNITIES



Non-Human Intelligence

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Artificial Intelligence



Our focus: **Human** intelligence because that's the intelligence we know...

Cognition: Perception, learning, reasoning, planning, and knowledge.

Deep learning is changing what we thought we could do, at least in perception and learning (with enough data).

Artificial Intelligence



Separate development --- **“non-human”**: Reasoning and planning. Similar qualitative and quantitative advances but **“under the radar.”**

Part of the world of software verification, program synthesis, and automating science and mathematical discovery.

Developments proceed **without** attempts to mimic human intelligence or even human intelligence capabilities.

Truly machine-focused (digital) on a task, e.g., “verify this code” or “synthesize this code” --- can use billions of inference steps --- or “synthesize an optimal plan with 1,000 steps.” (Near-optimal: 10,000+ steps.)

Example Consider a sequence of 1s and -1s, e.g.:

-1, 1, 1, -1, 1, 1, -1, 1, -1 ...
 1 2 3 4 5 6 7 8 9 ...
 2 4 6 8 ...
 3 6 9 ...

and look at the sum of sequences and subsequences:

$-1 + 1 = 0$

$-1 + 1 + 1 = 1$

$-1 + 1 + 1 + -1 = 0$

$-1 + 1 + 1 + -1 + 1 = 1$

$-1 + 1 + 1 + -1 + 1 + 1 = 2$ *

$-1 + 1 + 1 + -1 + 1 + 1 + -1 = 1$

$-1 + 1 + 1 + -1 + 1 + 1 + -1 + 1 = 2$

$-1 + 1 + 1 + -1 + 1 + 1 + -1 + 1 + -1 = 1$

and “skip by 1”

$1 + -1 = 0$

$1 + -1 + 1 = 1$

$1 + -1 + 1 + 1 = 2$

and “skip by 2”

$1 + 1 = 2$

$1 + 1 + -1 = 1$

* small inference step

We now know (2015): there exists a sequence of **1160** +1s and -1s such that sums of all subsequences *never* < -2 or > +2.

1160
elements
all sub-sums
stay between
-2 and +2

- + + - + - - + + - + + - + - - + + - + - - + - - +
+ - + - - + + - + + - + - + + - - + + - - + - - + + -
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40 x 29 pattern

So, we now know (2015): there exists a sequence of **1160** +1s and -1s such that sums of all subsequences *never* < -2 or $> +2$.

Result was obtained with a *general* propositional reasoning program (a Boolean Satisfiability or SAT solver). *Surprisingly*, the approach far outperformed specialized search methods written for the problem, including ones based on other known types of sequences. (A PolyMath project started in January 2010.)

Aside: A Taste of Problem Size

Consider a real world Boolean Satisfiability (SAT) problem, from formal verification.

The instance `bmc-ibm-6.cnf`, IBM LSU 1997:

`p cnf !`

`-1 7 0`

`-1 6 0`

`-1 5 0`

`-1 -4 0`

`-1 3 0`

`-1 2 0`

`-1 -8 0`

`-9 15 0`

`-9 14 0`

`-9 13 0`

`-9 -12 0`

`-9 11 0`

`-9 10 0`

`-9 -16 0`

`-17 23 0`

`-17 22 0`

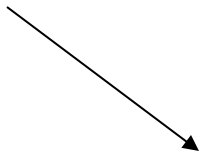
**i.e., ((not x_1) or x_7)
((not x_1) or x_6)
etc.**

***x_1, x_2, x_3 , etc. our Boolean variables
(set to True or False)***

Set x_1 to False ??

10 pages later:

1
185 -9 0
185 -1 0
177 169 161 153 145 137 129 121 113 105 97
89 81 73 65 57 49 41
33 25 17 9 1 -185 0
186 -187 0
186 -188 0
...



**I.e., (x_177 or x_169 or x_161 or x_153 ...
x_33 or x_25 or x_17 or x_9 or x_1 or (not x_185))**

clauses / constraints are getting more interesting...

Note x_1 ...

4000 pages later:

10236 -10050 0
10236 -10051 0
10236 -10235 0
10008 10009 10010 10011 10012 10013 10014
10015 10016 10017 10018 10019 10020 10021
10022 10023 10024 10025 10026 10027 10028
10029 10030 10031 10032 10033 10034 10035
10036 10037 10086 10087 10088 10089 10090
10091 10092 10093 10094 10095 10096 10097
10098 10099 10100 10101 10102 10103 10104
10105 10106 10107 10108 -55 -54 53 -52 -51 50
10047 10048 10049 10050 10051 10235 -10236 0
10237 -10008 0
10237 -10009 0
10237 -10010 0

...

Finally, 15,000 pages later:

-7 260 0
7 -260 0
1072 1070 0
-15 -14 -13 -12 -11 -10 0
-15 -14 -13 -12 -11 10 0
-15 -14 -13 -12 11 -10 0
-15 -14 -13 -12 11 10 0
-7 -6 -5 -4 -3 -2 0
-7 -6 -5 -4 -3 2 0
-7 -6 -5 -4 3 -2 0
-7 -6 -5 -4 3 2 0
185 0

Search space of truth assignments:

$$2^{50000} \approx 3.160699437 \cdot 10^{15051}$$

***Current SAT solvers solve this instance in
a few seconds!***

Back to sequences of +1/-1s

Encoding has variables for the sequence X_1, X_2, \dots, X_N

(we interpret True for +1 and False for -1)

but also e.g.

Proposition: “sum_of_first_2_terms_of_step_by_2_subseq = 2”

(for any given setting of $X_1 \dots X_N$ this is either True or False)

and statements of the form:

IF ((sum_of_first_2_terms_of_step_by_2_subseq = 2 == True)

AND (X_8 == False))

THEN

(sum_of_first_3_terms_of_step_by_2_subseq = 1 == True)

Encoding: 37,418 variables and 161,460 clauses / constraints.

Sequence found in about 1 hour (MacBook Air).

Perhaps SAT solver was “lucky” in finding the sequence?

But, remarkably, each sequence of 1161 or longer leads to a +3 (or -3) somewhere. (Erdos discrepancy conjecture)

Encoding: 37,462 variables and 161,644 clauses / constraints.

Proof of non-existence of discrepancy 2 sequence found in about 10 hour (MacBook Air).

Proof: 13 gigabytes and independently verified (50 line proof checking program). Proof is around a billion small inference steps.

E.g. Given $(A \text{ and } \sim B) \rightarrow C$

A

$\sim B$

conclude C.

Machine understands; humans: probably never. Still, we can be certain of the result because of the verifier.

Observations

- 1) Result different from earlier “computer math” results, such as the proof of the 4 color theorem, because here we don’t need to trust the theorem prover. Final proof (“certificate”) can be checked easily by anyone.
- 2) It’s **not a brute force search**. Earlier SAT solvers cannot find the proof. Specialized programs cannot find the proof.
Brute force proof is of order $2^{1161} = 3.13 \times 10^{349}$. Current solver finds complete proof with only around 1.2×10^{10} steps. Clever learning and reasoning enables a factor 10^{339} reduction in proof size (and discovery of the “short” proof).
- 3) In part inspired by discrepancy 2 result, Terence Tao proved just a few months ago the general Erdos conjecture (for any discrepancy). Deep and subtle math.
- 4) But, does not fully superseded the 1161 result for the discrepancy 2. Future math may build further on these types of computational results. (I.e. true, verifiable facts but not human accessible.)

Other examples



AlphaGo:

Core engine

Monte Carlo Tree Search (UCT, 2006)

Final boost: deep learning and reinforcement learning.

Search part *and insights* will likely remain beyond human understanding.

Planning: We can synthesize optimal plan sequences of 1000+ steps.

Changes the notion of a “program”

A planning-enabled robot will synthesize its plans on-the-fly given its current abilities. Quite different from current pre-programmed industrial robots.

Computational Complexity Hierarchy

EXP-complete:

games like Go, ...

PSPACE-complete:

QBF, *planning*, chess
(bounded), ...

#P-complete/hard:

#SAT, sampling,
probabilistic inference, ...

NP-complete:

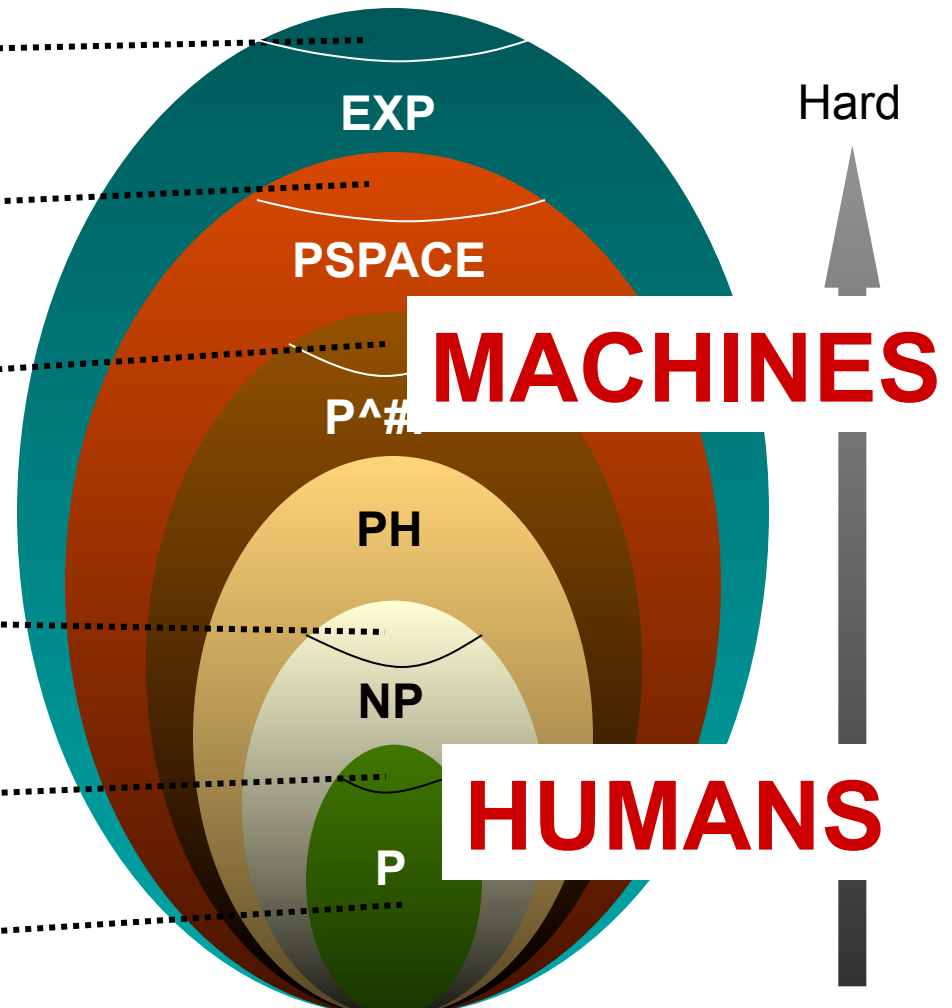
SAT, propositional
reasoning, scheduling,
graph coloring, puzzles, ...

P-complete:

circuit-value, ...

In P:

sorting, shortest path



**What are the consequences for human understanding
of machine intelligence?**

